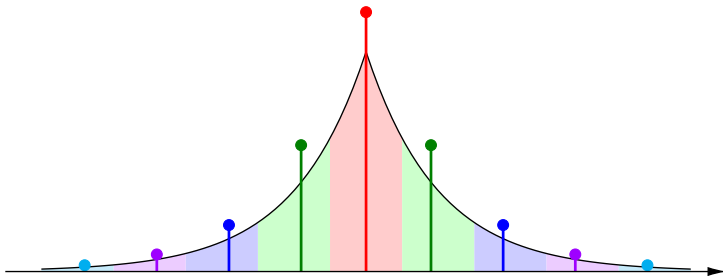


Quantization II



Entropy-Constrained Scalar Quantization (ECSQ)

- Lloyd quantizer: Minimize distortion for given number K of intervals
- Now: Consider quantizer design with variable-length coding of indices
- ➔ Average rate (without exploiting dependencies between quantization indices)

$$R = \sum_{\forall k} p_k \cdot \ell_k \geq H(S') = - \sum_{\forall k} p_k \log_2 p_k \quad (1)$$

with

$$p_k = \int_{u_k}^{u_{k+1}} f(s) ds \quad (2)$$

- ➔ Consider entropy instead of the rate of an actual code
- ➔ Average MSE distortion

$$D = \mathbb{E}\{d(S, Q(S))\} = \sum_{\forall k} \int_{u_k}^{u_{k+1}} (s - s'_k)^2 \cdot f(s) ds \quad (3)$$

Joint Minimization of Distortion and Rate/Entropy

- We look for solutions of constrained minimization problems

$$\min D \quad \text{subject to} \quad R \leq R_{\text{target}} \quad (4)$$

or equivalently

$$\min R \quad \text{subject to} \quad D \leq D_{\text{target}} \quad (5)$$

- ➔ Convert constrained optimization problem into unconstrained problem
- ➔ Minimized Lagrange function J (with Lagrange multiplier $\lambda > 0$)

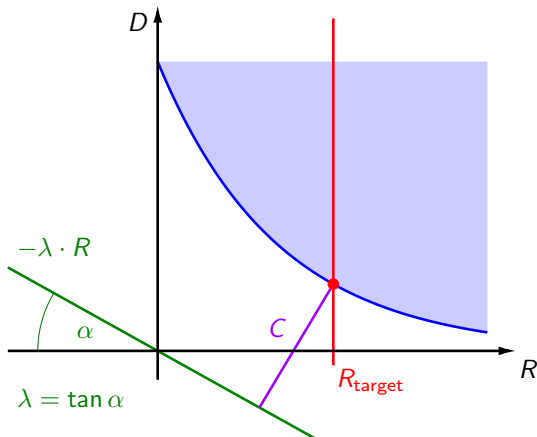
$$J = D + \lambda \cdot R \quad (6)$$

$$= \text{E}\{d(S, S')\} + \lambda \cdot \text{E}\{\ell(S')\} \quad (7)$$

$$= \text{E}\{d(S, S')\} + \lambda \cdot H(S') \quad (8)$$

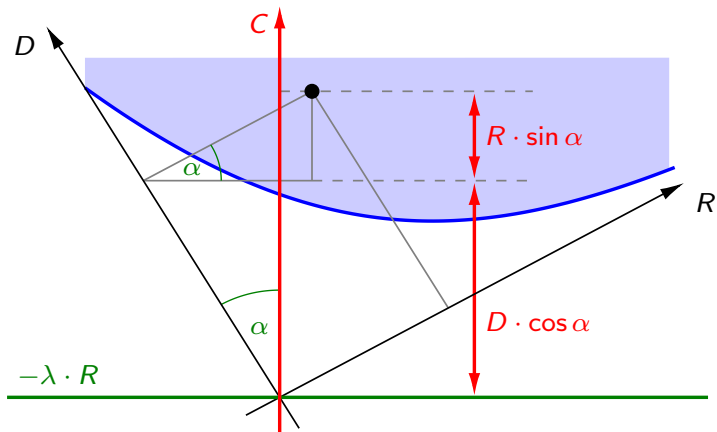
- Each Lagrange multiplier $\lambda > 0$ corresponds to a rate constraint R_{target} (or distortion constraint D_{target})

Lagrangian Optimization



- ➔ Points on convex hull: Minimize distance C to line $D = -\lambda \cdot R$
- ➔ Geometrical interpretation: Rotate coordinate system by angle α

Lagrangian Optimization



→ Minimize distance: $C = D \cdot \cos \alpha + R \cdot \sin \alpha$

→ Equivalent minimization: $J = D + \lambda \cdot R$ (note: $\lambda = \tan \alpha$)

Optimal Scalar Quantizer

- Minimize Lagrange function

$$\begin{aligned}
 J &= D + \lambda \cdot R \\
 &= \mathbb{E}\{d(S, S')\} + \lambda \cdot \mathbb{E}\{\ell(S')\} \\
 &= \mathbb{E}\{d(S, \beta(\alpha(S)))\} + \lambda \cdot \mathbb{E}\{\ell(\alpha(S))\} \\
 &= \sum_{\forall k} \int_{u_k}^{u_{k+1}} d(s, s'_k) f(s) ds + \lambda \cdot \sum_{\forall k} \ell_k \int_{u_k}^{u_{k+1}} f(s) ds \quad (9)
 \end{aligned}$$

- Lagrange parameter λ determines trade-off between rate and distortion
- Optimal design: Minimize $J = D + \lambda R$ with respect to
 - Reconstruction levels s'_k (decoder mapping β)
 - Decision thresholds u_k (encoding mapping α)
 - Codeword lengths ℓ_k (lossless mapping γ)

Optimal Decoder Mapping: Centroid Condition

$$J = \sum_{\forall k} \int_{u_k}^{u_{k+1}} d(s, s'_k) f(s) ds + \lambda \cdot \sum_{\forall k} \ell_k \int_{u_k}^{u_{k+1}} f(s) ds$$

1 Optimize reconstruction levels s'_k

for given decision thresholds u_k (and codeword lengths ℓ_k)

- Note: Rate term does not depend on reconstruction levels s'_k

→ **General centroid condition** (same as for Lloyd quantizer)

$$s'_k = \arg \min_{s'} \mathbb{E} \left\{ d(S, s') \mid S \in \mathcal{C}_k \right\} \quad (10)$$

→ **Centroid condition for MSE distortion**

$$s'_k = \mathbb{E} \{ S \mid S \in \mathcal{C}_k \} = \frac{1}{p_k} \int_{u_k}^{u_{k+1}} s f(s) ds = \frac{\int_{u_k}^{u_{k+1}} s f(s) ds}{\int_{u_k}^{u_{k+1}} f(s) ds} \quad (11)$$

Optimal Lossless Mapping: Entropy Condition

$$J = \sum_{\forall k} \int_{u_k}^{u_{k+1}} d(s, s'_k) f(s) ds + \lambda \cdot \sum_{\forall k} \ell_k \int_{u_k}^{u_{k+1}} f(s) ds$$

2 Optimize codeword lengths ℓ_k

for given decision thresholds u_k (and reconstruction levels s'_k)

- Note: Distortion term does not depend on codeword lengths ℓ_k
- ➔ Remember lossless coding: $R \geq H(S')$

$$\sum_{\forall k} p_k \cdot \ell_k \geq - \sum_{\forall k} p_k \cdot \log_2 p_k \quad (\text{equality iff } \ell_k = -\log_2 p_k) \quad (12)$$

- ➔ **Entropy condition** (neglecting inefficiency of actual entropy coding)

$$\ell_k = -\log_2 p_k = -\log_2 \left(\int_{u_k}^{u_{k+1}} f(s) ds \right) \quad (13)$$

Optimal Encoder Mapping

$$J = \sum_{\forall k} \int_{u_k}^{u_{k+1}} d(s, s'_k) f(s) ds + \lambda \cdot \sum_{\forall k} \ell_k \int_{u_k}^{u_{k+1}} f(s) ds$$

3 Optimize decision thresholds u_k

for given reconstruction levels s'_k and codeword lengths ℓ_k

→ Map each input value s to interval C_k that minimizes $d(s, s'_k) + \lambda \cdot \ell_k$

$$\alpha(s) = \arg \min_{\forall s'_k} d(s, s'_k) + \lambda \cdot \ell_k \quad (14)$$

■ Note: Each threshold u_k impacts only neighboring intervals C_{k-1} and C_k

→ **Modified nearest neighbor condition**

$$\boxed{d(u_k, s'_{k-1}) + \lambda \cdot \ell_{k-1} = d(u_k, s'_k) + \lambda \cdot \ell_k} \quad (15)$$

Optimal Encoder Mapping: MSE Distortion

- Optimal encoding mapping for MSE distortion

$$d(u_k, s'_{k-1}) + \lambda \cdot \ell_{k-1} = d(u_k, s'_k) + \lambda \cdot \ell_k$$

$$(u_k - s'_{k-1})^2 + \lambda \cdot \ell_{k-1} = (u_k - s'_k)^2 + \lambda \cdot \ell_k$$

$$u_k^2 - 2u_k s'_{k-1} + (s'_{k-1})^2 + \lambda \cdot \ell_{k-1} = u_k^2 - 2u_k s'_k + (s'_k)^2 + \lambda \cdot \ell_k$$

$$2u_k(s'_k - s'_{k-1}) = (s'_k)^2 - (s'_{k-1})^2 + \lambda(\ell_k - \ell_{k-1})$$

$$2u_k(s'_k - s'_{k-1}) = (s'_k - s'_{k-1})(s'_k + s'_{k-1}) + \lambda(\ell_k - \ell_{k-1})$$

- Optimal decision thresholds for MSE distortion

$$u_k = \frac{1}{2} (s'_{k-1} + s'_k) + \frac{\lambda}{2} \left(\frac{\ell_k - \ell_{k-1}}{s'_k - s'_{k-1}} \right) \quad (16)$$

- Threshold is shifted towards the reconstruction level with the longer codeword (less probable interval)

Optimal Entropy-Constrained Scalar Quantizer (ECSQ)

Necessary conditions for optimality

- Centroid condition (for reconstruction levels s'_k)

$$s'_k = \arg \min_{s' \in \mathbb{R}} \mathbb{E} \left\{ d(S, s') \mid S \in \mathcal{C}_k \right\} \quad (17)$$

- Entropy condition (for codeword lengths ℓ_k)

$$\ell_k = -\log_2 p_k = -\log_2 \int_{u_k}^{u_{k+1}} f(s) ds \quad (18)$$

- Modified nearest neighbor condition (for decision threshold u_k)

$$d(u_k, s'_{k-1}) + \lambda \cdot \ell_{k-1} = d(u_k, s'_k) + \lambda \cdot \ell_k \quad (19)$$

Optimal Entropy-Constrained Scalar Quantizer (ECSQ)

Optimality conditions for MSE distortion

- Centroid condition (for reconstruction levels s'_k)

$$s'_k = \frac{\int_{u_k}^{u_{k+1}} s f(s) ds}{\int_{u_k}^{u_{k+1}} f(s) ds} \quad (20)$$

- Entropy condition (for codeword lengths ℓ_k)

$$\ell_k = -\log_2 p_k \quad (21)$$

- Modified nearest neighbor condition (for decision threshold u_k)

$$u_k = \frac{1}{2} (s'_{k-1} + s'_k) + \frac{\lambda}{2} \left(\frac{\ell_k - \ell_{k-1}}{s'_k - s'_{k-1}} \right) \quad (22)$$

Design of optimal entropy-constrained scalar quantizers

- In general: Cannot be derived in closed form
- Iterative algorithm similar to Lloyd algorithm

Entropy-Constrained Lloyd Algorithm for Given Pdf (MSE)

- Given is:
- the marginal probability density function $f(s)$ of the source
 - a Lagrange multiplier $\lambda > 0$

Iterative quantizer design

- 1** Choose an initial set of reconstruction levels $\{s'_k\}$ and codeword lengths $\{\ell_k\}$
- 2** Update the decision thresholds $\{u_k\}$ according to

$$u_k = \frac{1}{2} (s'_{k-1} + s'_k) + \frac{\lambda}{2} \left(\frac{\ell_k - \ell_{k-1}}{s'_k - s'_{k-1}} \right)$$

- 3** Update the reconstruction levels $\{s'_k\}$ according to

$$s'_k = \frac{\int_{u_k}^{u_{k+1}} s f(s) ds}{\int_{u_k}^{u_{k+1}} f(s) ds}$$

- 4** Update the codeword lengths $\{\ell_k\}$ according to

$$\ell_k = -\log_2 p_k$$

- 5** Repeat the previous three steps until convergence

Entropy-Constrained Lloyd Algorithm for a Training Set

- Given is:
- a sufficiently large realization $\{s_n\}$ of the considered source
 - a Lagrange multiplier $\lambda > 0$

Iterative quantizer design

- 1 Choose an initial set of reconstruction levels $\{s'_k\}$ and codeword lengths $\{\ell_k\}$
- 2 Associate all samples of the training set $\{s_n\}$ with one of the quantization intervals \mathcal{C}_k according to

$$\alpha(s_n) = \arg \min_{\forall k} d(s_n, s'_k) + \lambda \cdot \ell_k$$

- 3 Update the reconstruction levels $\{s'_k\}$ according to

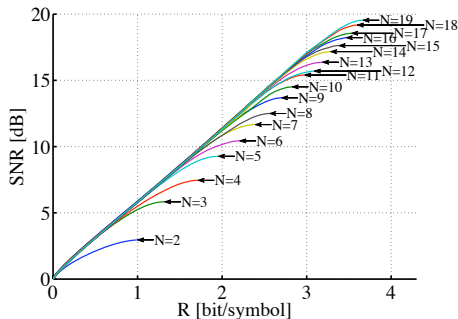
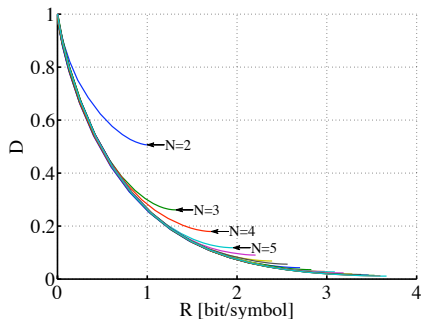
$$s'_k = \arg \min_{s' \in \mathbb{R}} \mathbb{E}\{d(S, s') \mid \alpha(S) = k\}$$

- 4 Update the codeword lengths $\{\ell_k\}$ according to

$$\ell_k = -\log_2 p_k$$

- 5 Repeat the previous three steps until convergence

Number of Initial Intervals for EC Lloyd Algorithm

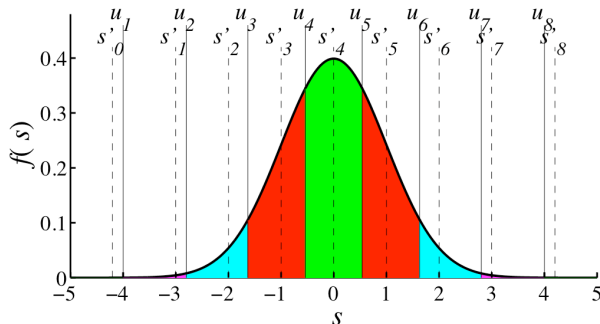


- Thresholds u_k are shifted towards less probable reconstruction levels
- Probability of less often chosen reconstruction levels is further reduced
- As a consequence, symbols get "removed" and the EC Lloyd algorithm can be initialized with more symbols than the final result

➔ Use large number of intervals in initialization

Entropy-Constrained Lloyd Algorithm for Gaussian Source

- Gaussian source with zero mean and unit variance
- Design optimal entropy-constrained quantizer with rate $R = 2$ bit/symbol



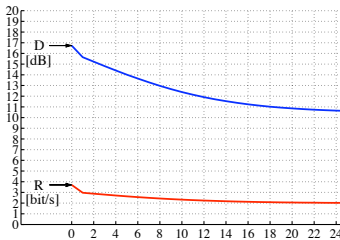
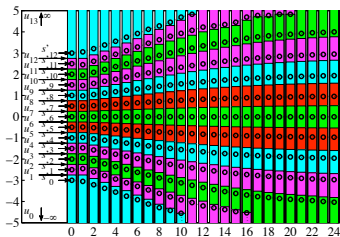
Comparison of MSE distortion (for $R = 2$ bits/sample)

- EC-Lloyd: $D = 0.09$ (SNR = 10.5 dB)
- Lloyd: $D = 0.12$ (SNR = 9.3 dB)
- SLB: $D = 0.0625$ (SNR = 12.0 dB)

Convergence of EC Lloyd Algorithm for Gaussian Source

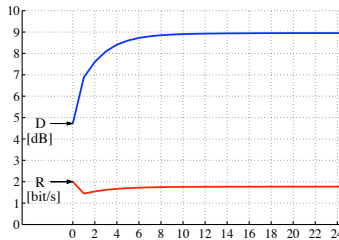
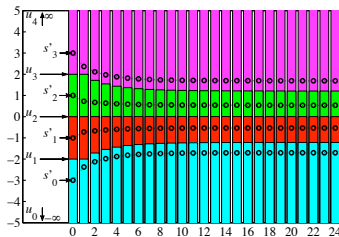
Initialization A:

$$s'_k = -3 + 0.5 \cdot k$$



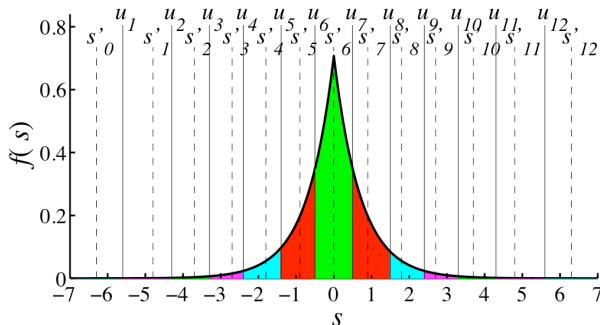
Initialization B:

$$s'_k = -3 + 2 \cdot k$$



Entropy-Constrained Lloyd Algorithm for Laplacian Source

- Laplacian source with zero mean and unit variance
- Design optimal entropy-constrained quantizer with rate $R = 2$ bit/symbol



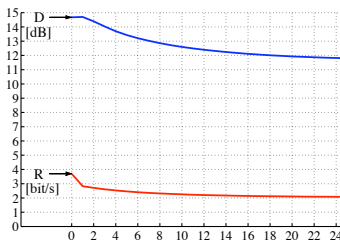
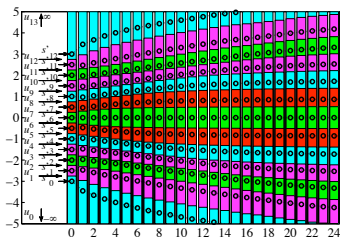
Comparison of MSE distortion (for $R = 2$ bits/sample)

- EC-Lloyd: $D = 0.07$ (SNR = 11.5 dB)
- Lloyd: $D = 0.18$ (SNR = 7.6 dB)
- SLB: $D = 0.054$ (SNR = 12.7 dB)

Convergence of EC Lloyd Algorithm for Laplacian Source

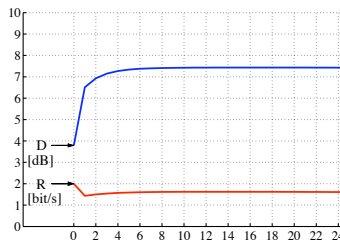
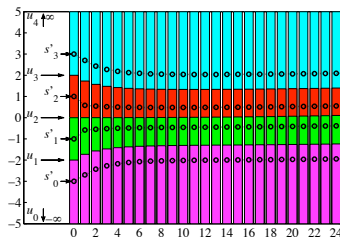
Initialization A:

$$s'_k = -3 + 0.5 \cdot k$$



Initialization B:

$$s'_k = -3 + 2 \cdot k$$



Remember: High-Rate Approximations

Approximations for High Rates / Small Distortions

- Marginal pdf $f(s)$ is (nearly) constant inside each quantization interval

$$f(s) \approx f(s'_k) \quad \forall s \in [u_k, u_{k+1}) \quad (23)$$

- ➔ Probability $p_k = P(S \in \mathcal{C}_k)$ is given by

$$p_k = \int_{u_k}^{u_{k+1}} f(s) ds = \Delta_k \cdot f(s'_k) \quad (24)$$

- ➔ Average distortion

$$D = \sum_{\forall k} \int_{u_k}^{u_{k+1}} d(s, s'_k) f(s) ds = \sum_{\forall k} f(s'_k) \int_{u_k}^{u_{k+1}} d(s, s'_k) ds \quad (25)$$

- ➔ Average MSE distortion for centroid quantizers

$$D = \frac{1}{12} \sum_{\forall k} p_k \Delta_k^2 \quad (26)$$

High-Rate Approximation: Optimal Quantizers with VLC

Average rate for optimal variable length coding (at high rates)

- Lossless source coding

$$R = H(S') = - \sum_{\forall k} p_k \log_2 p_k \quad (27)$$

- High rate approximation $p_k = f(s'_k) \Delta_k$

$$\begin{aligned} R &= - \sum_{\forall k} p_k (\log_2 f(s'_k) + \log_2 \Delta_k) \\ &= - \sum_{\forall k} f(s'_k) \log_2 f(s'_k) \Delta_k - \frac{1}{2} \sum_{\forall k} p_k \log_2 \Delta_k^2 \end{aligned} \quad (28)$$

- Asymptotic limit ($\Delta_k \rightarrow 0$) for first term

$$\begin{aligned} R &= - \int_{-\infty}^{\infty} f(s) \log_2 f(s) ds - \frac{1}{2} \sum_{\forall k} p_k \log_2 \Delta_k^2 \\ &= h(S) - \frac{1}{2} \sum_{\forall k} p_k \log_2 \Delta_k^2 \end{aligned} \quad (29)$$

High-Rate Approximation: Optimal Quantizers with VLC

- Will use: **Jensen's inequality for convex functions** $\psi(x)$

$$\sum_{\forall k} \alpha_k \psi(x_k) \geq \psi \left(\sum_{\forall k} \alpha_k x_k \right) \quad \text{for} \quad \sum_{\forall k} \alpha_k = 1 \quad (30)$$

with equality iff $x_k = \text{const}$

Average rate for optimal variable length coding (at high rates)

- Apply Jensen's inequality: Convex function $\psi(x) = -\log_2(x)$

$$\begin{aligned} R &= h(S) - \frac{1}{2} \sum_{\forall k} p_k \log_2 \Delta_k^2 \\ &\geq h(S) - \frac{1}{2} \log_2 \left(\sum_{\forall k} p_k \Delta_k^2 \right) \end{aligned} \quad (31)$$

with equality iff $\Delta_k = \text{const}$

High-Rate Approximation: Optimal Quantizers with VLC

Optimal scalar quantizers with variable length coding (at high rates)

- Remember: High-rate distortion approximation for centroid quantizers

$$D = \frac{1}{12} \sum_{\forall k} p_k \Delta_k^2$$

→ High-rate operational rate distortion function

$$R_V(D) = h(S) - \frac{1}{2} \log_2 \left(\sum_{\forall k} p_k \Delta_k^2 \right) = h(S) - \frac{1}{2} \log_2(12 D) \quad (32)$$

→ High-rate operational distortion rate function (Gish & Pierce, 1961)

$$D_V(R) = \frac{1}{12} 2^{2h(S)} 2^{-2R} \quad (33)$$

- For MSE distortion and high rates: Optimal scalar quantizer with variable length coding have uniform quantization step sizes

Comparison to Shannon Lower Bound

- High-rate approximation for optimal quantizer with variable-length codes

$$D_V(R) = \varepsilon_V^2 \cdot \sigma^2 \cdot 2^{-2R} \quad \text{with} \quad \varepsilon_V^2 = \frac{1}{12\sigma^2} \cdot 2^{2h(S)}$$

- Shannon lower bound: High-rate approximation of rate-distortion function

$$D_L(R) = \varepsilon_L^2 \cdot \sigma^2 \cdot 2^{-2R} \quad \text{with} \quad \varepsilon_L^2 = \frac{1}{2\pi e \sigma^2} \cdot 2^{2h(S)}$$

- Distortion increase (at same rate R) relative to Shannon lower bound

$$\frac{D_V(R)}{D_L(R)} = \frac{\varepsilon_V^2}{\varepsilon_L^2} = \frac{\pi e}{6} \approx 1.42 \quad (\text{Loss in SNR: 1.53 dB})$$

- Rate increase (at same distortion D) relative to Shannon lower bound

$$\begin{aligned} R_V(D) - R_L(D) &= \left(h(S) - \frac{1}{2} \log_2(12D) \right) - \left(h(S) - \frac{1}{2} \log_2(2\pi e D) \right) \\ &= \frac{1}{2} \log_2 \frac{\pi e}{6} \approx 0.2546 \quad (\approx 1/4 \text{ bit per sample}) \end{aligned}$$

Comparison of High-Rate Approximations

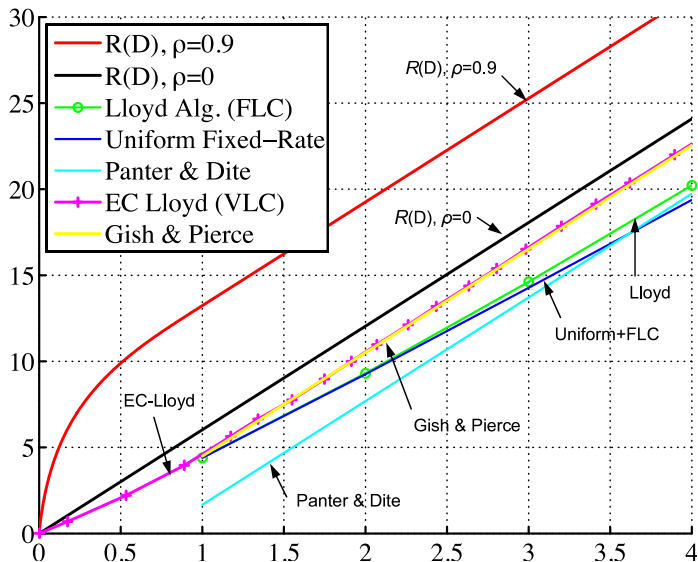
- Operational Distortion-rate function at high rates is given as

$$D_X(R) = \varepsilon_X^2 \cdot \sigma^2 \cdot 2^{-2R}$$

- Values of ε_X^2 depend on quantization method and pdf

<i>Method</i>	<i>Shannon Lower Bound (SLB)</i>	<i>Panter & Dite (Lloyd Quant. & FLC)</i>	<i>Gish & Pierce (ECSQ & VLC)</i>
Uniform pdf	$\frac{6}{\pi e} \approx 0.7$	1 (1.53 dB to SLB)	1 (1.53 dB to SLB)
Laplacian pdf	$\frac{e}{\pi} \approx 0.86$	$\frac{9}{2} = 4.5$ (7.1 dB to SLB)	$\frac{e^2}{6} \approx 1.23$ (1.53 dB to SLB)
Gaussian pdf	1	$\frac{\sqrt{3}\pi}{2} \approx 2.72$ (4.34 dB to SLB)	$\frac{\pi e}{6} \approx 1.42$ (1.53 dB to SLB)

Scalar Quantizers for Gaussian-Markov Sources



Summary

Lagrangian Optimization

- Re-formulate constrained optimization problems

$$\min_{\mathbf{p}} D(\mathbf{p}) \quad \text{subject to} \quad R(\mathbf{p}) \leq R_{\text{target}}$$

as unconstrained optimization problem (with $\lambda \geq 0$)

$$\min_{\mathbf{p}} D(\mathbf{p}) + \lambda \cdot R(\mathbf{p})$$

- Each solution \mathbf{p}_λ of the unconstrained problem does also represent a solution of the constrained problem with $R_{\text{target}} = R(\mathbf{p}_\lambda)$
- Lagrange multiplier λ determines trade-off between distortion and rate
- Geometrical interpretation:
 - Rotation of coordinate system by $\alpha = \arctan(\lambda)$
 - Conventional minimization in rotated space

Summary

Entropy-Constrained Scalar Quantizer (ECSQ)

- Optimal scalar quantizer (assuming optimal entropy coding)
- Three optimization criteria
 - Centroid condition (MSE): $s'_k = (\int_{u_k}^{u_{k+1}} s f(s) ds) / (\int_{u_k}^{u_{k+1}} f(s) ds)$
 - Entropy condition: $\ell_k = -\log_2 (\int_{u_k}^{u_{k+1}} f(s) ds)$
 - Mod. neighbor cond. (MSE): $u_k = \frac{1}{2}(s'_k + s'_{k-1}) + \frac{\lambda}{2} (\frac{\ell_k - \ell_{k-1}}{s'_k - s'_{k-1}})$
- Quantizer design: Iterative entropy-constrained Lloyd algorithm
 - Update decision thresholds u_k (for given s'_k and ℓ_k)
 - Update reconstruction levels s'_k and codeword lengths ℓ_k (for given u_k)
- May need large number ($K \rightarrow \infty$) of intervals for obtaining optimal quantizer
- Coding efficiency at high rates / small distortions
 - Distortion is factor 1.42 higher (1.53 dB) than SLB
 - Rate is roughly 0.25 bits per sample larger than SLB