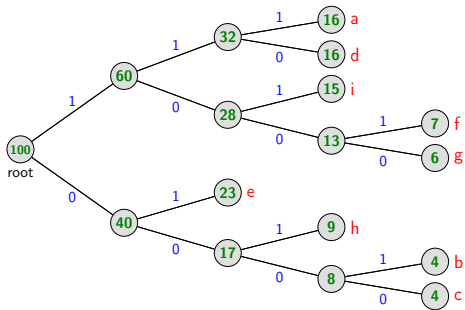


Lossless Coding

a_k	p_k	b_k
a	0.16	111
b	0.04	0001
c	0.04	0000
d	0.16	110
e	0.23	01
f	0.07	1001
g	0.06	1000
h	0.09	001
i	0.15	101



Morse Code (first version around 1837)

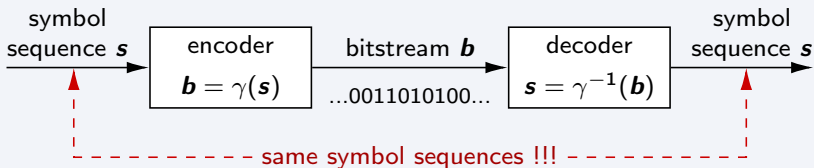
INTERNATIONAL MORSE CODE							
A	● —	N	— ●	1	● — — — —	.	● — ● — ● —
B	— ● ● ●	O	— — —	2	● ● — — —	,	— — ● ● — — —
C	— ● — ● ●	P	● — — ●	3	● ● ● — —	?	● ● — — ● ●
D	— ● ●	Q	— — — ● —	4	● ● ● ● —	!	● — — — — ●
E	●	R	● — ●	5	● ● ● ● ●		— ● — — ● — — —
F	● ● — ●	S	● ● ●	6	— ● ● ● ●	/	— ● ● — ●
G	— — — ●	T	—	7	— — ● ● ● ●	:	— — — — ● ● ●
H	● ● ● ●	U	● ● —	8	— — — — ● ●	;	— ● — — ● — — ●
I	● ●	V	● ● ● —	9	— — — — — ●	=	— ● ● ● —
J	● — — — —	W	● — — —	0	— — — — — —	+	● — ● — ●
K	— ● — —	X	— ● ● — —			-	— ● ● ● ● —
L	● — — ● ●	Y	— — ● — — —			_	● ● — — — ● —
M	— — —	Z	— — — ● ●			"	● — ● ● — — ●
						@	● — — — ● — ●

- Length of a dot is one unit
- Length of a dash is three units
- Space between parts of the same letter is one unit
- Space between letters is three units
- Space between words is seven units

Overview

Lossless Coding

- **Reversible mapping** of symbol sequences into bitstreams
- Each symbol must be taken from a countable alphabet (not necessarily finite)
- Original symbol sequence can be exactly recovered from bit sequence (in contrast to lossy coding)
- Also called **noiseless coding** or **entropy coding**



→ **Bit rate reduction only possible if source data have statistical properties that can be exploited for data compression**

Lossless Coding Framework

Framework

- Discrete random process $\mathbf{S} = \{S_n\}$ (discrete-time & discrete-amplitude)
- Consider messages $\mathbf{s}^{(L)}$ of L symbols

$$\mathbf{s}^{(L)} = \{s_0, s_1, s_2, \dots, s_{L-1}\}$$

- Messages represent finite-length realizations of random process \mathbf{S}

Lossless Coding

- Mapping of messages $\mathbf{s}^{(L)}$ onto bitstreams $\mathbf{b}^{(K)}$ with K bits

$$\mathbf{b}^{(K)} = \{b_0, b_1, b_2, \dots, b_{K-1}\} \quad \text{with} \quad \forall k, b_k \in \{0, 1\}$$

- Mapping γ from messages onto bitstreams is an invertible deterministic function

$$\mathbf{b}^{(K)} = \gamma(\mathbf{s}^{(L)}) \quad \iff \quad \mathbf{s}^{(L)} = \gamma^{-1}(\mathbf{b}^{(K)}) \quad (1)$$

Realization of Lossless Coding

Practical Realization

- Split messages $\mathbf{s}^{(L)}$ into fixed- or variable-length sub-sequences $\{\mathbf{s}_k^{(N_k)}\}$

$$\mathbf{s}_k^{(N_k)} = \{s_k, s_{k+1}, s_{k+2}, \dots, s_{k+N_k-1}\}$$

- Assign codewords $\mathbf{b}_k^{\ell_k}$ to sub-sequences $\mathbf{s}_k^{(N_k)}$
- Form bitstream $\mathbf{b}^{(K)}$ by concatenating the codewords

$$\mathbf{b}^{(K)} = \{\mathbf{b}_1^{(\ell_1)}, \mathbf{b}_2^{(\ell_2)}, \mathbf{b}_3^{(\ell_3)}, \dots\}$$

Lossless Code

- Encoder mapping:

$$\mathbf{b}_k^{(\ell_k)} = \gamma(\mathbf{s}_k^{(N_k)}) \quad (2)$$

- Decoder mapping:

$$\mathbf{s}_k^{(N_k)} = \gamma^{-1}(\mathbf{b}_k^{(\ell_k)}) \quad (3)$$

Example: Variable-Length Coding for Scalars

- Symbol alphabet: $\mathcal{A} = \{A, B, M, N\}$
- Lossless code: Codeword table

code A	
letter	codeword
A	00
B	01
M	10
N	11

code B	
letter	codeword
A	0
B	110
M	111
N	10

code C	
letter	codeword
A	0
B	010
M	100
N	10

- Example message: $s = \text{"BANANAMAN"}$
- ➔ Bitstream (code A): $b = \text{"010011001100100011"}$ (18 bits)
- ➔ Bitstream (code B): $b = \text{"1100100100111010"}$ (16 bits)
- ➔ Bitstream (code C): $b = \text{"0100100100100010"}$ (16 bits)

➔ **Goal: Minimize average codeword length**

$$\bar{\ell} = E\{\ell(S)\} = \sum_k p_k \cdot \ell_k \quad (4)$$

Example: Variable-Length Coding for Scalars

- Symbol alphabet: $\mathcal{A} = \{A, B, M, N\}$
- Lossless code: Codeword table

code A	
letter	codeword
A	00
B	01
M	10
N	11

code B	
letter	codeword
A	0
B	110
M	111
N	10

code C	
letter	codeword
A	0
B	010
M	100
N	10

Decoding

- Code A: $b = \text{"010011001100100011"}$ $\rightarrow s = \text{"BANANAMAN"}$
- Code B: $b = \text{"1100100100111010"}$ $\rightarrow s = \text{"BANANAMAN"}$
- Code C: $b = \text{"0100100100100010"}$ $\rightarrow s = \text{"B or AN ..."}$

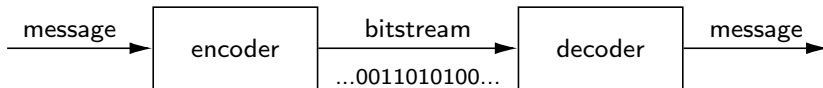
\rightarrow Necessary condition: Unique decodability:

Each bitstream uniquely represents a single message!

Outline

- Review: Probability, Random Variables, Random Processes
- Scalar Variable Length Codes
 - Prefix Codes
 - Unique Decodability
 - Entropy: Lower Limit for Average Codeword Length
 - Optimal Prefix Codes: Huffman Codes
 - Conditional Prefix Codes
- Variable-Length Codes for Vectors
 - Block Codes
 - V2V Codes
 - Entropy Rate
- Arithmetic Codes
 - Shannon-Fano-Elias Coding
 - Finite Precision: Arithmetic Coding

Mathematical Description of Source Coding



Transmission of new information to receiver

- Message is unknown by receiver
- ➔ Source can be modeled as a **random process**

Modeling of information sources as random processes

- ➔ Description using mathematical framework of probability theory
- ➔ Requires reasonable assumptions with respect to source of information
- ➔ Characterization of performance by probabilistic averages
- ➔ Basis for **mathematical theory of communication**

Probability Theory

- Branch of mathematics for describing and modeling random events
- Modern probability theory (the axiomatic definition of probability) introduced by KOLMOGOROV

Review:

- Probability
- Random variables
- Discrete random processes



ERGEBNISSE DER MATHEMATIK
UND IHRER GRENZGEBIETE
HERAUSGEGEBEN VON DER SCHIFFLEITUNG
DES
„ZENTRALBLATT FÜR MATHEMATIK“
ZWEITER BAND

GRUNDBEGRIFFE DER
WAHRSCHEINLICHS-
RECHNUNG

VON
A. KOLMOGOROFF



BERLIN
VERLAG VON JULIUS SPRINGER
1933

Random Experiment

Definition

- Any experiment with uncertain outcome ζ
- **Sample space** \mathcal{O} : Union of all possible outcomes ζ (**Certain event** \mathcal{O})
- **Event** \mathcal{A} : Union of zero or more possible outcomes ζ ($\mathcal{A} \subseteq \mathcal{O}$)

Examples

Roulette

- Sample space $\mathcal{O} = \{0, 32, 15, 19, 4, 21, \dots, 26\}$ (37 possible outcomes)
- Example event $\mathcal{A} = \{32\}$ (red 32)
- Example event $\mathcal{B} = \{15, 4, \dots, 26\}$ (any black number)

Voltage Measurement

- Sample space $\mathcal{O} = [0 \text{ V}; \infty \text{ V})$ (infinitely many possible outcomes)
- Example event $\mathcal{A} = [0 \text{ V}; 10 \text{ V})$ (not greater than 10 V)
- Example event $\mathcal{B} = (1 \text{ V}; 3 \text{ V})$ (between 1 V and 3 V)

Probability

Definition

- Measure $P(\mathcal{A})$ assigned to events \mathcal{A} of a random experiment
- Satisfies the following three axioms

Axioms

- 1 Probabilities are non-negative real numbers

$$P(\mathcal{A}) \geq 0, \quad \forall \mathcal{A} \subseteq \mathcal{O} \quad (5)$$

- 2 Certain event \mathcal{O} has a probability equal to 1

$$P(\mathcal{O}) = 1 \quad (6)$$

- 3 Probability of two disjoint events \mathcal{A} and \mathcal{B}

$$\mathcal{A} \cap \mathcal{B} = \emptyset \quad \implies \quad P(\mathcal{A} \cup \mathcal{B}) = P(\mathcal{A}) + P(\mathcal{B}) \quad (7)$$

Probability Estimation

Empirical Probability

- Repeatable random experiment
- Relative frequency of an event \mathcal{A} in N trials

$$\frac{N(\mathcal{A})}{N} = \frac{\text{number of trials in which } \mathcal{A} \text{ was observed}}{\text{number of total trials}} \quad (8)$$

- Empirical probability

$$P(\mathcal{A}) = \lim_{N \rightarrow \infty} \frac{N(\mathcal{A})}{N} \quad (9)$$

Practical Probability Estimation

- Repeat a random experiment N times
- Use the approximation

$$P(\mathcal{A}) \approx \frac{N(\mathcal{A})}{N}$$

→ Estimation quality depends on the size of N

Conditional Probability

Definition

- Probability of an event \mathcal{A} given that another event \mathcal{B} has occurred
- Referred to as **conditional probability** $P(\mathcal{A}|\mathcal{B})$
- Kolmogorov's definition

$$P(\mathcal{A}|\mathcal{B}) = \frac{P(\mathcal{A} \cap \mathcal{B})}{P(\mathcal{B})}, \quad \text{for } P(\mathcal{B}) > 0 \quad (10)$$

Bayes' Theorem

- For $P(\mathcal{A}) > 0$ and $P(\mathcal{B}) > 0$,

$$P(\mathcal{A}|\mathcal{B}) = P(\mathcal{B}|\mathcal{A}) \cdot \frac{P(\mathcal{A})}{P(\mathcal{B})} \quad (11)$$

- Can be directly derived from definition of conditional probability

$$P(\mathcal{A} \cap \mathcal{B}) = P(\mathcal{B} \cap \mathcal{A}) \implies P(\mathcal{A}|\mathcal{B}) \cdot P(\mathcal{B}) = P(\mathcal{B}|\mathcal{A}) \cdot P(\mathcal{A})$$

Independence of Events

Definition

- Two events \mathcal{A} and \mathcal{B} are said to be **independent** if and only if

$$P(\mathcal{A} \cap \mathcal{B}) = P(\mathcal{A}) \cdot P(\mathcal{B}) \quad (12)$$

Consequence

- Using reformulated definition of conditional probability

$$P(\mathcal{A} | \mathcal{B}) = \frac{P(\mathcal{A} \cap \mathcal{B})}{P(\mathcal{B})}, \quad \text{for } P(\mathcal{B}) > 0$$

- Two events \mathcal{A} and \mathcal{B} , with $P(\mathcal{B}) > 0$, are independent if and only if

$$P(\mathcal{A} | \mathcal{B}) = P(\mathcal{A}) \quad (13)$$

Random Variable

Definition

- A **random variable** S is a function of the sample space \mathcal{O} that assigns a real value $S(\zeta)$ to each outcome $\zeta \in \mathcal{O}$ of a random experiment

Examples for Random Variables

Roulette

- Random variable S : $S(\zeta) = \{\text{"number of pocket the ball lands"}\}$
- Random variable C : $C(\zeta) = \{0 : \text{"0"}, 1 : \text{"red"}, 2 : \text{"black"}\}$

Digital Camera

- Random variable V : $V(\zeta) = \{\text{"voltage (in V) of a photocell"}\}$
- Random variable Y : $Y(\zeta) = \{\text{"discrete sample value at output"}\}$

Note: A random variable may take ...

- a finite, a countably infinite, or an uncountable number of values

Cumulative Distribution Function

Definition

- **Cumulative distribution function** (cdf) of a random variable S

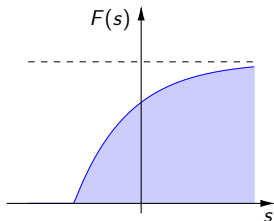
$$F_S(s) = P(S \leq s) = P(\{\zeta : S(\zeta) \leq s\}) \quad (14)$$

- $F_S(s)$ is also referred to as **distribution** of a random variable S

Properties

- $F_S(s)$ is non-decreasing
- $F_S(-\infty) = 0$
- $F_S(\infty) = 1$
- $P(a < S \leq b) = F_S(b) - F_S(a)$

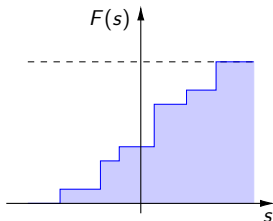
Examples: Cumulative Distribution Functions



Continuous function

- Random variable S can take all values inside one or more non-zero intervals

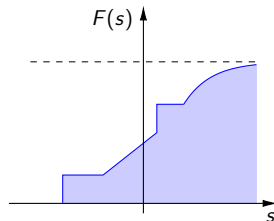
→ **Continuous random variable**



Staircase function

- Random variable S can only take a countable number of values

→ **Discrete random variable**



Mixed type

- Random variable S can take all values inside one or more intervals and a countable number of additional values

Joint Cumulative Distribution Function

Multiple Random Variables

- Joint cdf / joint distribution of two random variables X and Y

$$F_{XY}(x, y) = P(X \leq x, Y \leq y) \quad (15)$$

- Straightforward extension to $N > 2$ random variables

Random Vectors

- Random vector \mathbf{S} : Vector of multiple random variables
- N -dimensional random vector: $\mathbf{S} = \mathbf{S}^{(N)} = \{S_0, S_1, \dots, S_{N-1}\}$
- N -dimensional joint cdf / joint distribution

$$F_{\mathbf{S}}(\mathbf{s}) = P(\mathbf{S} \leq \mathbf{s}) = P(S_0 \leq s_0, S_1 \leq s_1, \dots, S_{N-1} \leq s_{N-1}) \quad (16)$$

- Joint cdf / joint distribution of two random vectors \mathbf{X} and \mathbf{Y}

$$F_{\mathbf{XY}}(\mathbf{x}, \mathbf{y}) = P(\mathbf{X} \leq \mathbf{x}, \mathbf{Y} \leq \mathbf{y}) \quad (17)$$

Conditional Cumulative Distribution Function

Conditional Cdfs / Conditional Distribution

- Conditional cdf of a random variable S given event \mathcal{B} , with $P(\mathcal{B}) > 0$

$$F_{S|\mathcal{B}}(s|\mathcal{B}) = P(S \leq s | \mathcal{B}) = \frac{P(\{S \leq s\} \cap \mathcal{B})}{P(\mathcal{B})} \quad (18)$$

- Conditional cdf of a random variable X given another random variable Y

$$F_{X|Y}(x|y) = P(X \leq x | Y \leq y) = \frac{P(X \leq x, Y \leq y)}{P(Y \leq y)} = \frac{F_{XY}(x, y)}{F_Y(y)} \quad (19)$$

Random Vectors

- Conditional cdf of a random vector \mathbf{X} given another random vector \mathbf{Y}

$$F_{\mathbf{X}|\mathbf{Y}}(\mathbf{x}|\mathbf{y}) = P(\mathbf{X} \leq \mathbf{x} | \mathbf{Y} \leq \mathbf{y}) = \frac{P(\mathbf{X} \leq \mathbf{x}, \mathbf{Y} \leq \mathbf{y})}{P(\mathbf{Y} \leq \mathbf{y})} = \frac{F_{\mathbf{X}\mathbf{Y}}(\mathbf{x}, \mathbf{y})}{F_{\mathbf{Y}}(\mathbf{y})} \quad (20)$$

Discrete Random Variables: Probability Mass Function

Discrete Random Variables

- A random variable S is called a **discrete random variable**, if and only if its cdf $F_S(s)$ is a staircase function
- S can only take values of a countable alphabet

$$\mathcal{A}_S = \{a_0, a_1, \dots\} \quad (21)$$

Definition: Probability Mass Function of a Discrete Random Variable S

- Probability mass function (pmf) is defined on alphabet \mathcal{A}_S

$$p_S(a) = P(\{\zeta \in \mathcal{O} : S(\zeta) = a\}) \quad (22)$$

- Property

$$\sum_{\forall a \in \mathcal{A}_S} p_S(a) = P(\mathcal{O}) = 1 \quad (23)$$

- Simplified notation

$$p_k = p_S(a_k) \quad (24)$$

Cdf for Discrete Random Variables

Cumulative Distribution Function

- Cdf for discrete random variables S

$$F_S(s) = \sum_{\forall a \in \mathcal{A}_S : a \leq s} p_S(a) \quad (25)$$

Cumulative Mass Function

- Discrete version of cdf
- Cumulative mass function (cmf) defined on alphabet $\mathcal{A}_S = \{a_0, a_1, \dots\}$

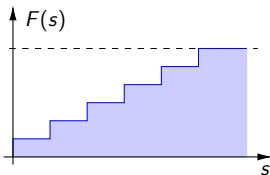
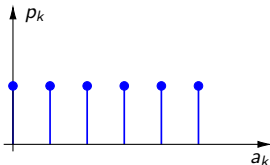
$$c_k = c(a_k) = \sum_{\forall a \in \mathcal{A}_S : a \leq a_k} p_S(a) \quad (26)$$

Examples for Discrete Distributions

Uniform

$$p_k = \frac{1}{M}$$

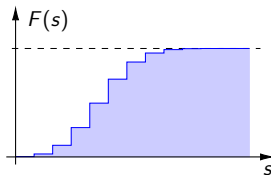
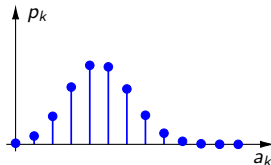
$$(0 \leq k < M)$$



Binomial

$$p_k = \binom{n}{k} p^k (1-p)^{n-k}$$

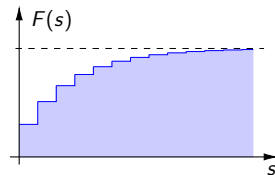
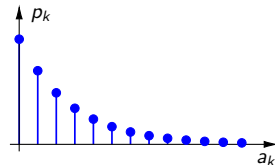
$$(0 \leq k \leq n)$$



Geometric

$$p_k = (1-p)^k p$$

$$(k \geq 0)$$



Joint Probability Mass Function

Multiple Random Variables

- Joint pmf of two discrete random variables X and Y

$$p_{XY}(x, y) = P(X = x, Y = y) \quad (27)$$

- Straightforward extension to $N > 2$ random variables

Random Vectors

- Consider random vector $\mathbf{S} = \mathbf{S}^{(N)} = \{S_0, S_1, \dots, S_{N-1}\}$

- N -dimensional joint pmf

$$p_{\mathbf{S}}(\mathbf{s}) = P(\mathbf{S} = \mathbf{s}) \quad (28)$$

- Joint pmf of two random vectors \mathbf{X} and \mathbf{Y}

$$p_{\mathbf{X}\mathbf{Y}}(\mathbf{x}, \mathbf{y}) = P(\mathbf{X} = \mathbf{x}, \mathbf{Y} = \mathbf{y}) \quad (29)$$

Conditional Probability Mass Function

Conditional Pmf

- Conditional pmf of a random variable S given event \mathcal{B} , with $P(\mathcal{B}) > 0$

$$p_{S|\mathcal{B}}(a|\mathcal{B}) = P(S = a|\mathcal{B}) = \frac{P(\{S = a\} \cap \mathcal{B})}{P(\mathcal{B})} \quad (30)$$

- Conditional pmf of a random variable X given another random variable Y

$$p_{X|Y}(x|y) = P(X = x|Y = y) = \frac{P(X = x, Y = y)}{P(Y = y)} = \frac{p_{XY}(x, y)}{p_Y(y)} \quad (31)$$

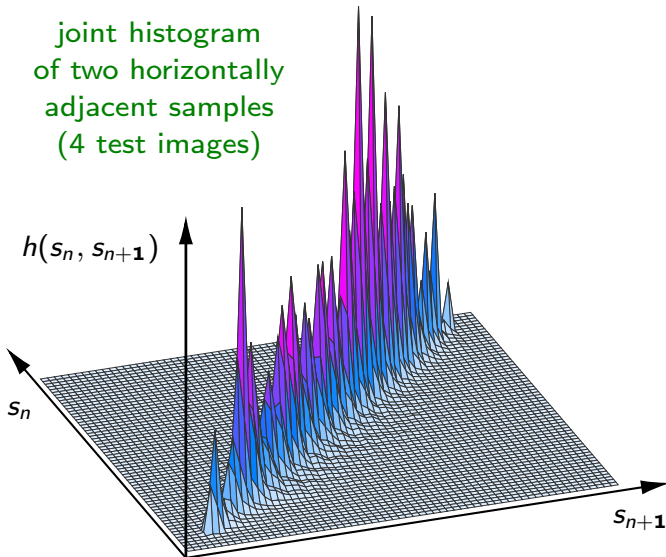
Random Vectors

- Conditional pmf of a random vector \mathbf{X} given another random vector \mathbf{Y}

$$p_{\mathbf{X}|\mathbf{Y}}(\mathbf{x}|\mathbf{y}) = \frac{p_{\mathbf{X}\mathbf{Y}}(\mathbf{x}, \mathbf{y})}{p_{\mathbf{Y}}(\mathbf{y})} \quad (32)$$

Example for a Joint Pmf (Empirical Estimation)

joint histogram
of two horizontally
adjacent samples
(4 test images)



Expected Value / Expectation

Definition

- Expected value of a function $g(S)$ of a random variable S
- **Discrete random variable** S with alphabet \mathcal{A}

$$E\{g(S)\} = \sum_{\forall a \in \mathcal{A}} g(a) p_S(a) \quad (33)$$

Multiple Random Variables

- Expected value of function $g(X, Y)$ of two random variable X and Y

$$E\{g(X, Y)\} = \sum_{x,y} g(x, y) p_{XY}(x, y) \quad (34)$$

- Straightforward extension to more than two random variables

Conditional Expectations

Given an Event

- Expected value of function $g(S)$ given event \mathcal{B} , with $P(\mathcal{B}) > 0$

$$E\{g(S) | \mathcal{B}\} = \sum_a g(a) p_{S|\mathcal{B}}(a | \mathcal{B}) \quad (35)$$

- Expected value of $g(X)$ given a value y for another random variable Y

$$E\{g(X) | y\} = E\{g(X) | Y = y\} = \sum_x g(x) p_{X|Y}(x, y) \quad (36)$$

→ $E\{g(X) | y\}$ is a deterministic function of y

Conditional Expectations

Given another Random Variable

- Expected value of $g(X)$ given another random variable Y

$$E\{g(X) | Y\} = \sum_x g(x) p_{X|Y}(x, Y) \quad (37)$$

→ $E\{g(X) | Y\}$ is another random variable

Properties of Expected Values

Important Properties

- Linearity

$$E\{ aX + bY \} = aE\{ X \} + bE\{ Y \} \quad (38)$$

- For independent random variables X and Y

$$E\{ XY \} = E\{ X \} E\{ Y \} \quad (39)$$

- Iterative expectation rule

$$E\{ E\{ g(X) | Y \} \} = E\{ g(X) \} \quad (40)$$

Proofs: *Homework*.

Important Expected Values

Mean, Variance, Covariance

- **Mean** μ_S and **variance** σ_S^2 of a random variable S

$$\mu_S = E\{S\} \quad \text{and} \quad \sigma_S^2 = E\{(S - \mu_S)^2\} \quad (41)$$

- **Covariance** σ_{XY}^2 of two random variables X and Y , and **correlation coefficient** ϕ_{XY} between two random variables X and Y

$$\sigma_{XY}^2 = \text{cov}(X, Y) = E\{(X - \mu_X)(Y - \mu_Y)\} \quad (42)$$

$$\phi_{XY} = \frac{\sigma_{XY}^2}{\sigma_X \sigma_Y} \quad (43)$$

Average Codeword Length

- Average codeword length for given lossless code

$$\bar{\ell} = E\{\ell(S)\} = \sum_{a \in \mathcal{A}} p(a) \cdot \ell(a) \quad (44)$$

Random Processes

Discrete-Time Random Process

- Series of random experiments at time instants t_n , with $n = 0, 1, 2, \dots$
- For each experiment: Random variable $S_n = S(t_n)$
- **Random process**: Series of random variables

$$\mathbf{S} = \{S_n\}$$

Characterization of Statistical Properties

- Consider N -dimensional random vector $\mathbf{S}_k^{(N)} = \{S_k, S_{k+1}, \dots, S_{k+N-1}\}$
- N -th order joint cdf

$$F_k^{(N)}(\mathbf{s}) = \mathrm{P}\left(\mathbf{S}_k^{(N)} \leq \mathbf{s}\right) = \mathrm{P}(S_k \leq s_0, S_{k+1} \leq s_1, \dots, S_{k+N-1} \leq s_{N-1}) \quad (45)$$

Discrete Random Processes

Discrete Random Processes

- Random variables $\{S_n\}$ are of discrete type
- Each random variable S_n has an alphabet \mathcal{A}_n
- N -th order joint pmf

$$p_k^{(N)}(\mathbf{a}) = \text{P}(S_k = a_0, S_{k+1} = a_1, \dots, S_{k+N-1} = a_{N-1}) \quad (46)$$

with

$$F_k^{(N)}(\mathbf{s}) = \sum_{\forall a_0 \in \mathcal{A}_0: a_0 \leq s_0, \dots} p_k^{(N)}(a_0, a_1, \dots, a_{N-1}) \quad (47)$$

Stationary Process

Stationary Random Processes

- Statistical properties are **invariant to a shift in time**
- ➔ Joint cdf, pmf do not depend on start index k

$$\forall k, i : \quad F_k^{(N)}(\mathbf{s}) = F_i^{(N)}(\mathbf{s}) \quad (48)$$

$$p_k^{(N)}(\mathbf{a}) = p_i^{(N)}(\mathbf{a}) \quad (49)$$

- In this course: Only consider stationary processes
- ➔ Simplify notations
 - N -th order cdf: $F_N(\mathbf{s})$
 - N -th order pmf: $p_N(\mathbf{a})$

Memoryless and IID Processes

Memoryless Random Process

- Random process $\mathbf{S} = \{S_n\}$ for which all random variables S_n are independent of each other

Independent and Identically Distributed (IID) Random Process

- **Stationary** and **memoryless**
- With $F_S(s)$ and $p_S(a)$ being the cdf and pmf for a single random variable $S = S_n$ (called marginal cdf and pmf), we have

$$F_N(\mathbf{s}) = \prod_{k=0}^{N-1} F_S(s_k), \quad p_N(\mathbf{a}) = \prod_{k=0}^{N-1} p_S(a_k) \quad (50)$$

Markov Processes

Random Process with Markov Property

- Markov Property: Future outcomes do not depend on past outcomes, but only on the present outcome

$$P(S_n \leq s_n | S_{n-1} \leq s_{n-1}, \dots) = P(S_n \leq s_n | S_{n-1} \leq s_{n-1}) \quad (51)$$

Stationary Markov Process

- Stationary random process with Markov Property
- Statistical properties completely specified by 1-st order conditional cdf

$$F(s_n | s_{n-1}) = P(S_n \leq s_n | S_{n-1} \leq s_{n-1}) \quad (52)$$

- With marginal cdf $F_S(s)$, the N -th order joint cdf is given by

$$F_N(\mathbf{s}) = F_S(s_0) \prod_{k=1}^{N-1} F(s_k | s_{k-1}) \quad (53)$$

Stationary Discrete Markov Processes

Markov Property for Discrete Random Processes

- Statistical properties completely specified by 1-st order conditional pmf

$$p(a_n | a_{n-1}) = P(S_n = a_n | S_{n-1} = a_{n-1}) \quad (54)$$

- With marginal pmf $p_S(a)$, the N -th order joint pmf is given by

$$p_N(\mathbf{a}) = p_S(a_0) \prod_{k=1}^{N-1} p(a_k | a_{k-1}) \quad (55)$$

Example: Stationary Discrete Markov Process

Specified by conditional pmf $p(a_n | a_{n-1})$

a	a_0	a_1	a_2
$p(a a_0)$	0.90	0.05	0.05
$p(a a_1)$	0.15	0.80	0.05
$p(a a_2)$	0.25	0.15	0.60

Homework:

Determine marginal pmf
 $p_S(a) = P(S = a)$

Summary

Lossless Coding

- Reversible mapping: Symbol sequence \iff bitstream
- Necessary condition: Unique decodability
- Optimization criterion: Minimize average codeword length

Probability

- Axiomatic definition, empirical probability
- Conditional probability and independence

Discrete Random Variables

- Cumulative distribution function (cdf): Staircase function
- Probability mass function (pmf)
- Expected values: Mean, variance, covariance, average codeword length

Discrete Random Processes

- Stationary, memoryless, iid, Markov processes